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Abstract

In a previous report by Grune a.o. [1] the practical use of two algorithms is illustrated. The first algorithm determines which relations, such as "contains" or "may follow", may hold between the terminal productions of notions and/or symbols of a given context-free grammar, the second one verifies the LL(1)-ness of context-free grammars. In this report the ALGOL 60 programs that perform the corresponding calculations are given, together with some small examples.

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1. Introduction

In [1], some methods are explained that are employed by our group in the construction of an ALGOL 68 compiler. Amongst others, an algorithm is given to determine which relations, such as "contains" or "may follow", may hold between the terminal productions of notions and/or the symbols of a given context-free grammar, and an algorithm to verify the LL(1)-ness of context-free grammars. In fact, both these algorithms are well-known and have already been investigated by many others (see, e.g., [2,3]). Whereas in [1] the emphasis lies on the illustration of the practical use of those algorithms, in this report the actual programs (written in ALGOL 60) are given that perform the corresponding calculations. They are termed: "Relations concerning a cf-grammar" and "LL(1)-checker", respectively.

The input of both programs is a context-free grammar of the following form: first, all terminal symbols of the grammar are given, each one separated from the next by a semicolon, the last one followed by a point. Second, the production rules of the grammar are given, one rule for each non-terminal. Those rules should be given according to the following grammar:

```
rule: left hand side, colon, alternatives, point.
left hand side: notion.
alternatives: alternative option; alternative option, semicolon,
alternatives.
alternative option: alternative;.
alternative: member; member, comma, alternative.
member: notion; open, notion list, close.
notion list: notion; notion, comma, notion list.
```

Whenever a list of notions is placed between the symbols "open" "(" and "close" (")", this makes the list optional. Layout characters and comment (placed between the braces "[" and "]") are skipped, except for layout characters within notions, in which case 1 space is honoured to increase readability of the (quite bulky) output.

The author would like to thank D. Grune and L. Meertens for their valuable suggestions.

1.1 Reading the input

The input is read by a simple parser, which replaces each notion by some number for further use. A correct grammar is parsed correctly and no incorrect grammar passes unnoticed. As a notion may consist of only 27 distinct characters (26 letters and the space symbol), a 32-radix number system is used (i.e., 5 bits per symbol). On our machine (an EL X8 of Philips-Electrologica) one word contains 27 bits, so 5 symbols are packed per word. On other machines, some adaptations may be necessary: `max int` (= maximum integral number), part of the routine `next symbol`, which performs the packing, and the routine `get notion at`, which does the unpacking.

Reading the input is done through the external procedure `resymbol`. Each time `resymbol` is called, the character code of the next symbol of the input stream is delivered. Our character code is such that the letters a through z are represented by a contiguous set of integers (from the value of letter a up to letter a + 25). This property is used by the routines `next symbol` (packing), `letter` and `get notion at` (unpacking). All other character codes used are represented by variables with suggestive identifiers, initialized at the beginning to the proper values.

1.2 Other machine-dependent features

Both programs make use of the following output procedures:

procedure `printtext` (`s`); string `s`;

The string `s` is printed.

procedure `absfixt` (`n`, `m`, `a`); value `n`, `m`, `a`; integer `n`, `m`; real `a`;

For `m = 0`, the integral value of `a` is printed in `n` positions without a sign.

procedure `nlcr`;

A new line is generated.

procedure carriage (n); value n; integer n;

This procedure generates n new lines.

procedure newpage;

A new page is generated.

procedure space (n); value n; integer n;

This procedure generates n spaces.

procedure prsym (n); value n; integer n;

The printing equivalent of resymbol. The character corresponding to the value of n, is printed.

integer procedure printpos;

An informative procedure, delivering the current position on the line. The value delivered is 0 in case the line is empty. The positions on the line are numbered from 0 up to "line width". Thus, it indicates the first free position.

integer procedure linenum;

The value of the function designator linenum is the number of the line we are currently printing on. Thus, its value is 1 after a new page is given.

In addition, the programs make use of four other machine-dependent procedures:

procedure exit;

A call of this procedure terminates the program.

integer procedure available;

The value of this function designator is the number of words in memory that is still available.

integer procedure stringsymbol (i,s); value i; integer i; string s;

The character code of the i-th symbol of the string s is delivered, where i counts from 0.

procedure fix (n, m, x, a); value n, m, x; integer n, m; real a;
integer array a;

For $m = 0$, the digits of the number in x are placed in consecutive elements of a, starting with a [2] (a [1] contains the sign of x), up to a [n + 2] (which always contains a space).

On the output of the program "Relations concerning a cf-grammar", the linewidth is adjusted to 70. This can be changed by adjusting the integer variable line width.

Some remarks should be given about transferring those programs to other machines. On our machine, booleans are packed, i.e., the word length is 27 bits, so 27 booleans are packed in one word. But still, a grammar with say 250 distinct notions could not be run without more (in the initialization part of both programs, the integer variable max notion is given a value indicating an upperbound of the number of distinct notions according to which the size of the arrays is determined; the remaining space is used for storing the notions and the grammar). In one of the programs, as listed here, ten huge boolean matrices are used. They are, however, not necessary in parallel, so some of them may share memory, and can be swapped in and out whenever necessary. However, on a machine which does not pack booleans, it will, even then, probably be impossible to digest such a grammar. In this case the user has to do the packing himself and rewrite the routines transitive closure, mul 2 and mul 3 as well, and adjust several others.

2. Relations concerning a cf-grammar

The relations, dealt with in this program, are "may contain", "may be contained in", "may begin with", "may be the begin of", "may end with", "may be the end of" (, all of which are transitive,) and "may follow" and "may precede". A short description of the outcome of this program may suffice here; a more extensive description, together with an elaborate example, can be found in [1].

The first relation given is "may contain". We say α "WITHIN" β if β occurs in the right hand side of the production rule for α . "may contain" may then be defined as: may contain = WITHIN⁺, where the ⁺ indicates the transitive closure of the considered relation.

We say α "FIRST" β if β occurs as the first member of some direct production of α . "may begin with" is then given by: may begin with = FIRST⁺. Similarly, we define "LAST", so that: may end with = LAST⁺. We say α "FOLLOW" β if, in some production rule the succession of members $\beta, \gamma_1, \dots, \gamma_n$, α occurs, $n \geq 0$ and γ_i produces empty ($1 \leq i \leq n$). "may follow" can then be computed as follows: may follow = FIRST⁺ FOLLOW LAST⁺, that is, α may follow β if there exist γ and δ such that γ FOLLOW δ , $\gamma = \alpha$ or γ may begin with α , and $\delta = \beta$ or δ may end with β . (The * indicates the reflexive transitive closure.) The other relations given are simply the transpose of one of the relations treated above.

For each relation R a list is given of entries for each of the notions, where an entry of the form "notion - $s_1; \dots; s_n$." indicates that $\{s | \text{notion } R s\} = \{s_1, \dots, s_n\}$. An entry of the form "notion₁, ..., notion_m - $s_1; \dots; s_n$." is an abbreviation of m entries of the form "notion_i - $s_1; \dots; s_n$.", $1 \leq i \leq m$.

In each entry, the part " $s_1; \dots; s_n$." is split up into two parts, the first one of the form " $s_1; \dots; s_j$.", containing only terminals (and preceded by the string (terminals:)), the second one of the form " $s_{j+1}; \dots; s_n$.", containing only non-terminals (and preceded by "(non-terminals:)"). If one of those parts is empty, the string preceding it is also omitted.

On the output, the entries of each relation are numbered, this number, together with a two-letter code preceding it, being a reference for the alphabetical listing of the notions given at the end.

In an entry, one or more of the s_i are marked with an asterisk. If, for a relation R , some s_j is unmarked, this implies that there is a marked s_i such that $s_i Q s_j$, where Q is R (is "may begin with", is "may end with") if R is a transitive relation (is "may follow", is "may precede"). E.g., for the relation FIRST^+ , the subset which is FIRST is marked. Also, one or more of the notion_i may be marked with a plus, indicating that this relation R is reflexive with respect to notion_i , i.e., $\text{notion}_i R \text{notion}_i$ holds.

In section 2.1. the ALGOL-60-text of the program is given, in section 2.2. a small example of a grammar and the resulting output is presented.

2.1 The program

```

begin comment relations concerning a cf-grammar, september 1973;
  integer end of file, eof, space char, tab char, nlcr char,
    sub char, bus char, stock, nil, comma char, comma,
    point char, point, colon char, colon, semicolon char,
    semicolon, open char, open, close char, close, empty,
    list pointer, max int, first notion, new defined,
    head of tree, top, error count, sym, number, numb,
    notion, max notion, start notion, harmless, minus char,
    n of adm cells, upperbound, i, star, symbols, letter a,
    space repr, bits per word, line width, plus char,
    zero char;
  integer array t32[0 : 4];

initialization part:
  end of file:= eof:= -4096; space char:= 93; tab char:= 118;
  nlcr char:= 119; sub char:= 100; bus char:= 101; nil:= -1;
  comma char:= comma:= 87; point char:= 88; point:= -5;
  colon char:= colon:= 90; semicolon char:= 91; semicolon:= -4;
  open char:= 98; open:= -2; close char:= 99; close:= -3; empty:= 0;
  max int:= 67 108 863; first notion:= 1; new defined:= -1;
  t32[0]:= 1 048 576; t32[1]:= 32 768; t32[2]:= 1 024; t32[3]:= 32;
  t32[4]:= 1; list pointer:= 1; star:= 66; error count:= 0;
  harmless:= -1; n of adm cells:=10; letter a:= 10; zero char:= 0;
  minus char:= 65; space repr:= 27; bits per word:= 27;
  line width:= 70; plus char:= 64;
  max notion:= 300;
  upperbound:= top:= available - max notion - 10 × max notion ×
    max notion / bits per word - 2000;
  printtext(↑upperbound array list:↑); absfixt(5, 0, upperbound);
  nlcr; nlcr;

begin integer array list[1 : upperbound],
  notion link[1 : max notion];
  comment 'notion link' contains pointers to notions, stored in
    'list', as follows:
    left
    right
    7 places for the index
    number(also used for the index)
  pointer: text 1
    :
    text last - max int,
  starting from 1 up to 'list pointer'.
  the essential part of the grammar is stored in 'list',
  from 'upperbound' to 'top' (i.e. backwards);
  boolean array first, first star, last, last star, follow,
    within, within star, aux1, aux2, aux3
    [1 : max notion, 1 : max notion],
    possibly empty[1 : max notion];

```

comment reading department;

integer procedure char;

begin integer i, j;

repeat: i:= char:= resymbol;

if i ≠ end of file then prsym(i);

if i = tab char v i = nlcr char then goto repeat else

if i = space char then

begin j:= resymbol;

for i:= j while i = space char v i = tab char v

 i = nlcr char v i = sub char do

begin if i = sub char then

begin for i:= sub char, j while j ≠ bus char ^

 j ≠ end of file do

begin prsym(i); j:= resymbol end;

if j = bus char then

begin prsym(j); j:= resymbol end

else error(←premature end of file→)

end

else

begin prsym(i); j:= resymbol end

end;

if j ≠ end of file then prsym(j);

if letter(j) then stock:= j else char:= j

end

else if i = sub char then

 skip comment:

begin i:= char:= resymbol; prsym(i);

if i = bus char then goto repeat;

if i ≠ end of file then goto skip comment

end

end char;

procedure next symbol;

comment this procedure yields a symbol in "sym", and,
 if it is a notion, yields its number in "numb";

begin

again: if sym = end of file then else

if stock = nil then sym:= char else

begin sym:= stock; stock:= nil end;

if sym = space char then goto again;

if sym = comma char then sym:= comma else

if sym = point char then sym:= point else

if sym = colon char then sym:= colon else

if sym = semicolon char then sym:= semicolon else

if sym = open char then sym:= open else

if sym = close char then sym:= close else

if sym = end of file then sym:= eof else

if letter(sym) then

begin integer j, aux, sm;

 number:= number + 1;

```

for j:= nil, nil, empty, empty, empty, empty, empty,
    empty, empty, number do store(j);
aux:= j:= 0; start notion:= list pointer;
for sm:= sym, sm while letter(sm) do
begin if j = 5 then
    begin store(aux); aux:= j:= 0 end store in word;
    aux:= t32[j] × (if sm = space char then space repr
    else sm - letter a + 1) + aux; j:= j + 1;
    if stock = nil then sm:= char else
    begin sm:= stock; stock:= nil end
end pack symbols;
store( - max int + aux); stock:= sm;
numb:= in tree(start notion, head of tree);
if numb = new defined then
begin numb:= new(number); notion link[number]:= - start
    notion
end not yet defined
else
if defined(numb) then numb:= new(numb);
sym:= notion
end read and store a notion
else
begin error({illegal character; skipped}); goto again end
end next symbol;

```

```

integer procedure in tree(start, last ref);
value start, last ref; integer start, last ref;
comment this procedure considers the notion at "start". if this
    notion is in the tree, it yields its number and removes
    the notion, otherwise, it adds the notion to the tree
    and yields "new defined";
begin integer node, comp;
    boolean found;

```

```

integer procedure compare(one, two); value one, two;
integer one, two;
comment compare:= sign(one '-' two);
begin integer o, t, so, st, i;
    o:= list[one]; t:= list[two];
    for i:= 1 while o = t ^ o > 0 do
    begin one:= one + 1; o:= list[one]; two:= two + 1;
        t:= list[two]
    end;
    comment discriminating item found, now analyze it;
    if o < 0 then
    begin o:= o + max int; so:= - 1 end
    else so:= 1;
    if t < 0 then
    begin t:= t + max int; st:= - 1 end
    else st:= 1;
    compare:= sign(if o = t then so - st else o - t)
end compare;

```

```

    found:= false;
    comment search tree;
    for node:= list[last ref] while node ≠ nil ^ not found do
    begin comp:= compare(node, start);
        if comp = 0 then found:= true else
            last ref:= node - n of adm cells +
                (if comp < 0 then 1 else 0)
    end node;
    comment analyze result: ;
    if found then
    begin listpointer:= start - n of adm cells;
        number:= number - 1; comment item removed;
        in tree:= list[node - 1]
    end list[head - 1] contains the number of the notion
    else
    begin list[last ref]:= start; in tree:= new defined
    end item added
end in tree;

boolean procedure letter(sym); value sym; integer sym;
letter:= letter a ≤ sym ^ sym < letter a + 26 ∨ sym = space char;

procedure store(sym); value sym; integer sym;
if list pointer > top then
begin carriage(2); printtext({space exhausted}); exit end
else
begin list[list pointer]:= sym; list pointer:= list pointer + 1
end store;

procedure store on top(sym); value sym; integer sym;
if top < list pointer then
begin carriage(2); printtext({space exhausted}); exit end
else
begin list[top]:= sym; top:= top - 1 end store on top;

integer procedure new(number); value number; integer number;
new:= - number;

boolean procedure is new(number); value number;
integer number;
is new:= number < 0;

procedure define(number); value number; integer number;
notion link[number]:= abs(notion link[number]);

boolean procedure defined(number); value number;
integer number;
defined:= notion link[number] > 0;

```

comment the syntax of the input is given by:

```

grammar: terminals, rules, eof.

terminals: notion, point|
           notion, semicolon, terminals.

rules: rule, (rules).
rule: notion, colon, tail, point.
tail: (alternative), (semicolon, tail).
alternative: member, (comma, member).
member: notion|
        open, notion list, close.
notion list: notion, (comma, notion list).

```

the next part reads in such a grammar, puts its structure in the back end of the array "list", and puts the notions in a tree, stored at the front end of "list", obviously, the | stands for ;

```

procedure grammar;
begin number:= 0; sym:= stock:= nil; head of tree:= list pointer;
      store(nil); next symbol; terminals; rules
end grammar;

```

```

procedure terminals;
terminals:
begin req notion(false);
      if is new(num) then define(new(num)) else
        error({terminal occurs twice});
      if is(semicolon, false) then goto terminals else
      if is(point, false) then symbols:= number - 1 else
      if is(eof, false) then else
        begin error({error in terminals}); next symbol;
        goto terminals
      end
end terminals;

```

```

procedure rules;
rules:
begin req notion(true);
  if is new(num) then define(new(num)) else
    error({notion already defined});
  if is(colon, false) then else error({colon missing});
tail: if is(point, true) then
  begin if is(eof, false) then else goto rules end
  else if is(semicolon, true) then goto tail else
    if member then
rest list:
  begin if is(comma, false) then
    begin if ¬ member then error({alternative wrong});
      goto rest list
    end
    else goto tail
  end
  else if is(eof, false) then
    begin store on top(point); error({premature end of file});
    end
    else
      begin error({incorrect rule}); next symbol; goto tail end
end rules;

boolean procedure member;
if is(notion, true) then member:= true else
if is(open, true) then
begin member:= true;
rest member: req notion(true);
  if is(comma, false) then goto rest member else
  if is(close, true) then else error({option wrong})
end
else member:= false;

procedure req notion(copy); value copy; boolean copy;
if ¬ is(notion, copy) then
begin error({notion missing}); numb:= harmless end req notion;

boolean procedure is(s, copy); value s, copy; integer s;
boolean copy;
if s = sym then
begin is:= true; if copy then
  store on top(if sym = notion then abs(num) else sym);
  next symbol
end
else is:= false;

procedure error(s); string s;
begin integer pos;
  pos:= printpos; error count:= error count + 1; carriage(2);
  printtext({error: }); printtext(s); carriage(3); space(pos)
end error;

```

```

comment compute "possibly empty", "within", "within star",
    "first", "first star", "last", "last star" and "follow";

procedure check for empty productions;
comment check which notions produce empty. the procedure continues
    until there are no more changes;
begin integer aux, notion, el, i;
    boolean changed, any;
    any := false;
    for i := 1 step 1 until number do
        possibly empty[i] := false;
    check for empty productions: changed := false;
    for aux := upperbound, aux - 1 while aux > top do
        begin notion := list[aux];
            if possibly empty[notion] then
                begin for el := list[aux] while el ≠ point do
                    aux := aux - 1
                end was already empty
            else
                begin
                    next: aux := aux - 1; el := list[aux]; if el = open then
                        begin for el := list[aux] while el ≠ close do
                            aux := aux - 1; goto next
                        end optional part
                    else if el = semicolon ∨ el = point then
                        begin for el := list[aux] while el ≠ point do
                            aux := aux - 1;
                            possibly empty[notion] := changed := true
                        end empty production found
                    else if possibly empty[el] then goto next else
                        for el := list[aux] while el ≠ point do
                            if el = semicolon then goto next else
                                aux := aux - 1
                            end production rule
                        end grammar;
                    if changed then goto check for empty productions;
                    if any then
                        begin printtext(⌘the following rules may produce empty⌘);
                            nlcr;
                            for i := 1 step 1 until number do
                                if possibly empty[i] then
                                    begin nlcr; print notion(i) end
                                end
                            else printtext(⌘no rule produces empty⌘); newpage
                        end check for empty productions;
                    end
                end
            end
        end
    end

```

```

procedure may contain;
comment the relation "directly contains" is stored in "within",
    its transitive closure in "within star";
begin integer i, j, aux, el, lhs;
    for i:= 1 step 1 until number do
    for j:= 1 step 1 until number do
    within[i, j]:= within star[i, j]:= false;
    for aux:= upperbound, aux - 1 while aux > top do
    begin lhs:= list[aux]; aux:= aux - 1;
        for el:= list[aux] while el ≠ point do
        begin if el > 0 then
            within[lhs, el]:= within star[lhs, el]:= true;
            aux:= aux - 1
        end rule
    end grammar;
    transitive closure(within star)
end may contain;

procedure may begin with;
comment this procedure determines the relation "may begin with".
    the result is stored in 'first' (directly beginning with),
    its transitive closure in 'first star';
begin integer i, j, el, aux, notion;
    boolean optional, errors;
    errors:= false;
    for i:= 1 step 1 until number do
    for j:= 1 step 1 until number do
    first[i, j]:= first star[i, j]:= false;
    for aux:= upperbound, aux - 1 while aux > top do
    begin notion:= list[aux]; aux:= aux - 1; optional:= false;
        for el:= list[aux] while el ≠ point do
        begin aux:= aux - 1;
            if el = open then optional:= true else
            if el = close then optional:= false else
            if el ≠ semicolon then
            begin first[notion, el]:= first star[notion, el]:=
                true;
                if possibly empty[el] then else
                if optional then
                begin for el:= list[aux] while el ≠ close
                    do aux:= aux - 1
                end skip rest of option
                else
                for el:= list[aux] while el ≠ semicolon ^
                    el ≠ point do aux:= aux - 1
            end
        end
    end production rule
end grammar;
    transitive closure(first star)
end may begin with;

```

```

procedure may end with;
comment this procedure determines the relation "may end with";
begin integer i, j, n, p, aux, notion, el;
    boolean optional;
    for i:= 1 step 1 until number do
    for j:= 1 step 1 until number do
        last[i, j]:= last star[i, j]:= false;
    for aux:= upperbound, aux while aux > top do
        begin notion:= list[aux]; n:= aux; optional:= false;
            for el:= list[aux] while el ≠ point, el do
                aux:= aux - 1; p:= aux + 2;
                for el:= list[p] while p < n do
                    begin p:= p + 1;
                        if el = open then optional:= false else
                            if el = close then optional:= true else
                                if el ≠ semicolon then
                                    begin last[notion, el]:= last star[notion, el]:=
                                        true;
                                        if possibly empty[el] then else
                                            if optional then
                                                begin for el:= list[p] while el ≠ open do
                                                    p:= p + 1
                                                end skip rest of option
                                            else
                                                for el:= list[p] while el ≠ semicolon ^ p < n
                                                    do p:= p + 1
                                            end
                                        end production rule
                                    end grammar;
                                transitive closure(last star)
                            end may end with;
                    end
                end
            end
        transitive closure(r); boolean array r;
        comment warshall algorithm, see "grammar handling tools", p48;
        begin integer i, j, k;
            for i:= 1 step 1 until number do
                for j:= symbols + 1 step 1 until number do
                    if r[j, i] then
                        begin for k:= 1 step 1 until number do
                            if r[i, k] then r[j, k]:= true
                        end
                    end
                end
            end transitive closure;

```

```

procedure mul2(f, v, r); boolean array f, v, r;
comment v transposed x f is computed;
begin integer a, b, c;
  for a:= 1 step 1 until number do
    for b:= 1 step 1 until number do r[a, b]:= false;
    for a:= 1 step 1 until number do
      for c:= 1 step 1 until number do
        if f[c, a] then
          begin for b:= 1 step 1 until number do
            if v[c, b] then r[a, b]:= true
          end
        end
      end
    end
  end mul2;

```

```

procedure mul3(f, v, l, r); boolean array f, v, l, r;
comment "multiplication" of 3 boolean arrays, the result is
  stored in "r". see "grammar handling tools", p. 49;
begin integer a, b, c, d;
  for a:= 1 step 1 until number do
    for b:= 1 step 1 until number do r[a, b]:= false;
    for c:= 1 step 1 until number do
      for d:= 1 step 1 until number do
        if v[c, d] then
          begin for a:= 1 step 1 until number do
            if f[c, a] then
              begin for b:= 1 step 1 until number do
                if l[d, b] then r[a, b]:= true
              end b, a
            end d, c
          end
        end
      end
    end
  end mul3;

```

```

procedure transpose(r) in:(r transp);
boolean array r, r transp;
begin integer i, j;
  for i:= 1 step 1 until number do
    for j:= 1 step 1 until number do r transp[i, j]:=r[j, i]
  end transpose;

```

```

procedure follow within;
comment this procedure determines the successions of notions
  within the production rules;
begin integer i, j, stackpointer, lwb stack, upb stack, el,
  previous, aux;
  integer array stack[1 : 100];
  comment the upperbound 100 may be erroneous, a good one is
    the maximum number of notions in an alternative,
    but this will very likely be less than 100;
  for i:= 1 step 1 until number do
    for j:= 1 step 1 until number do follow[i, j]:= false;

```

```

for aux:= upperbound, aux - 1 while aux > top do
for el:= semicolon, list[aux] while el ≠ point do
begin aux:= aux - 1;
  if el = open then
begin if previous ≠ nil then
  begin stackpointer:= stackpointer + 1;
    stack[stackpointer]:= previous
  end;
  upb stack:= stackpointer + 1
end
else if el = close then lwb stack:= upb stack:= 1
else if el = semicolon then
begin stackpointer:= 0; lwb stack:= upb stack:= 1;
  previous:= nil
end
else
begin for i:= lwb stack step 1 until stackpointer
do follow[el, stack[i]]:= true;
  if previous ≠ nil then
begin follow[el, previous]:= true;
  if possibly empty[el] then
begin stackpointer:= stackpointer + 1;
  stack[stackpointer]:= previous
end
else
begin stackpointer:= upb stack - 1;
  lwb stack:= upb stack
end
end;
  previous:= el
end notion
end rule and grammar
end follow within;

comment output department;

procedure print list of notions;
begin integer i;
  printtext(⟨list of notions, each notion preceded by its number⟩);
  nlcr; printtext(⟨and, if not defined by an asterisk⟩); nlcr;
  for i:= 1 step 1 until number do
begin nlcr; if ¬ defined(i) then
  begin prsym(star); define(i); error count:= error count+1
end
else prsym(spacechar);
  absfixt(4, 0, i); print notion(i)
end
end print list of notions;

```

```
procedure print notion(number); value number; integer number;
print notion at(notion link[number]);
```

```
procedure print notion at(index); value index; integer index;
begin integer i, n;
    integer array a[1 : 100];
    comment this upperbound is rather arbitrary, notions of
        length > 100 are considered rare;
    get notion at(index) of length:(n) in:(a);
    for i:= 1 step 1 until n do prsym(a[i])
end print notion at;
```

```
procedure get notion at(index) of length:(n) in:(a);
value index; integer index, n; integer array a;
begin integer el, j, k;
    n:= 0;
    for el:= list[index] while el > 0, el + max int do
    begin for k:= 0, k + 1 while k < 5 ^ el > 0 do
        begin j:= el div t32[k]; el:= el - j × t32[k]; n:= n + 1;
            a[n]:= if j = space repr then space char else
                j + letter a - 1
        end;
        index:= index + 1
    end
end get notion;
```

```
procedure print relation(r, marked, heading, index,
                        kind of relation);
value index; integer index; boolean array r, marked;
string heading, kind of relation;
begin integer i, notion, count, aux;
    boolean in lhs, first terminal, first nonterminal;
    boolean array printed, union marked[1 : number];
```

```
    procedure newline;
    begin nlcr; if linenumber = 1 then
        begin printtext(heading); nlcr; nlcr end
    end heading;
```

```
    procedure print notion(number); value number;
    integer number;
    begin integer n, i;
        integer array a[1 : 100];
        comment see remark at "print notion at";
        get notion at(notion link[number]) of length:(n) in:(a);
        if if printpos < 8 then false else
            printpos + n > line width then
            begin newline; space(if in lhs then 5 else 8) end;
            for i:= 1 step 1 until n do prsym(a[i])
        end print notion;
```

```

for i:= 1 step 1 until number do printed[i]:= false;
newpage; printtext(heading); nlcr; nlcr; count:= 0;
for notion:= 1 step 1 until number do
  if  $\neg$  printed[notion] then
    begin aux:= count:= count + 1; in lhs:= true;
      space(if count < 10 then 2 else
        if count < 100 then 1 else 0);
      printtext(kind of relation);
      if aux > 99 then
        begin i:= aux div 100; prsym(i + zero char);
          aux:= aux - i  $\times$  100
        end;
      if aux > 9 then
        begin i:= aux div 10; prsym(i + zero char);
          aux:= aux - i  $\times$  10
        end;
      prsym(aux + zero char); prsym(space char);
      print notion(notion);
      list[notion link[notion] - index]:= count;
      if r[notion, notion] then prsym(plus char);
      for i:= 1 step 1 until number do
        union marked[i]:= marked[notion, i];
      for i:= notion + 1 step 1 until number do
        if  $\neg$  printed[i] then
          begin boolean the same;
            the same:= true; aux:= 0;
            for aux:= aux + 1 while aux < number ^ the same
              do the same:= r[notion, aux] = r[i, aux];
            if the same then
              begin prsym(comma char); prsym(space char);
                print notion(i);
                list[notion link[i] - index]:= count;
                if r[i, i] then prsym(plus char);
                for aux:= 1 step 1 until number do
                  if marked[i,aux] then union marked[aux]:=true;
                printed[i]:= true
              end the same rhs, so taken together
            end grouping;
            prsym(minus char); newline; in lhs:= false;
            first terminal:= first nonterminal:= true;

```

```

    for i:= 1 step 1 until number do
    if r[notion, i] then
    begin if first terminal ^ i < symbols then
        begin printtext(† (terminals:) †);
            first terminal:= false
        end first terminal
    else if first nonterminal ^ i > symbols then
    begin if ¬ first terminal then
        begin prsym(point char); newline end close ;
        printtext(† (nonterminals:) †);
            first nonterminal:= false
        end first nonterminal
    else
        begin prsym(semicolon char); prsym(space char) end;
        print notion(i);
        if union marked[i] then prsym(star)
        end print notion i;
        prsym(point char); newline; newline
    end notion
end print relation;

procedure print index(head); value head; integer head;
begin integer i, j;
    boolean first;
    integer array a[1 : 5];
    if list[head - n of adm cells] ≠ nil then
    print index(list[head - n of adm cells]);
    print notion at(head);
    if printpos < 28 then space(28 - printpos) else
    begin nlcr; space(28) end;
    first:= true;
    for j:= 0 step 1 until 7 do
    begin if ¬ first then
        begin prsym(comma char); prsym(space char) end
        else first:= false;
        for i:= 4 × j, i + 1 do
        prsym(stringsymbol(i, †mc, ci, bw, bo, ew, eo, mf, mp†));
        fix(3, 0, list[head + j - 8], a);
        for i:= 1 step 1 until 5 do
        if a[i] < letter a then prsym(a[i] + zero char)
        end;
        nlcr; if list[head - n of adm cells + 1] ≠ nil then
        print index(list[head - n of adm cells + 1])
    end print index;

```

main program:

```

grammar;
newpage;
print list of notions;
if error count > 0 then
begin carriage(3); absfixt(3, 0, error count);
  printtext(4 errors in grammar); exit
end;
newpage;
check for empty productions;
may contain;
print relation(within star, within, {may contain (mc): ab - a; b.},
  8, {mc});
transpose(within star) in:(aux1);
transpose(within) in:(aux2);
print relation(aux1, aux2, {may be contained in (ci): a, b - ab.},
  7, {ci});

may begin with;
print relation(first star, first, {may begin with (bw): ab - a.},
  6, {bw});
transpose(first star) in:(aux1);
transpose(first) in:(aux2);
print relation(aux1, aux2, {may be the begin of (bo): a - ab.},
  5, {bo});

may end with;
print relation(last star, last, {may end with (ew): ab - b.},
  4, {ew});
transpose(last star) in:(aux1);
transpose(last) in:(aux2);
print relation(aux1, aux2, {may be the end of (eo): b - ab.},
  3, {eo});

follow within;
for i:= 1 step 1 until number do
first star[i, i]:= last star[i, i]:= true;
mul3(first star, follow, last star, aux1);
mul2(first star, follow, aux2);
print relation(aux1, aux2, {may follow (mf): b - a.}, 2, {mf});
transpose(aux1) in:(aux3);
transpose(follow) in:(aux1);
mul2(last star, aux1, aux2);
print relation(aux3, aux2, {may precede (mp): a - b.}, 1, {mp});
newpage;
print index(list[head of tree])

```

end
end

2.2 An example

UPPERBOUND ARRAY LIST: 20431

[INPUT:]

[TERMINAL SYMBOLS]

COLON; POINT; SEMICOLON; COMMA; OPEN; CLOSE; NOTION; EOF.

GRAMMAR: TERMINALS, RULES, EOF.

TERMINALS: NOTION, REST TERMINALS.

REST TERMINALS: POINT; SEMICOLON, TERMINALS.

RULES: RULE, (RULES).

RULE: LEFT HAND SIDE, COLON, ALTERNATIVES, POINT.

LEFT HAND SIDE: NOTION.

ALTERNATIVES: (ALTERNATIVE), (SEMICOLON, ALTERNATIVES).

ALTERNATIVE: MEMBER, (COMMA, ALTERNATIVE).

MEMBER: NOTION; OPEN, NOTION LIST, CLOSE.

NOTION LIST: NOTION, (COMMA, NOTION LIST).

MAY CONTAIN (MC): AB - A; B.

MC1 COLON, POINT, SEMICOLON, COMMA, OPEN, CLOSE, NOTION, EOF-

MC2 GRAMMAR-

(TERMINALS:) COLON; POINT; SEMICOLON; COMMA; OPEN; CLOSE; NOTION;
EOF*.

(NONTERMINALS:) TERMINALS*; RULES*; REST TERMINALS; RULE;
LEFT HAND SIDE; ALTERNATIVES; ALTERNATIVE; MEMBER; NOTION LIST.

MC3 TERMINALS+, REST TERMINALS+-

(TERMINALS:) POINT*; SEMICOLON*; NOTION*.
(NONTERMINALS:) TERMINALS*; REST TERMINALS*.

MC4 RULES+-

(TERMINALS:) COLON; POINT; SEMICOLON; COMMA; OPEN; CLOSE; NOTION.
(NONTERMINALS:) RULES*; RULE*; LEFT HAND SIDE; ALTERNATIVES;
ALTERNATIVE; MEMBER; NOTION LIST.

MC5 RULE-

(TERMINALS:) COLON*; POINT*; SEMICOLON; COMMA; OPEN; CLOSE; NOTION.
(NONTERMINALS;) LEFT HAND SIDE*; ALTERNATIVES*; ALTERNATIVE;
MEMBER; NOTION LIST.

MC6 LEFT HAND SIDE-

(TERMINALS:) NOTION*.

MC7 ALTERNATIVES+-

(TERMINALS:) SEMICOLON*; COMMA; OPEN; CLOSE; NOTION.
(NONTERMINALS;) ALTERNATIVES*; ALTERNATIVE*; MEMBER; NOTION LIST.

MC8 ALTERNATIVE+-

(TERMINALS:) COMMA*; OPEN; CLOSE; NOTION.
(NONTERMINALS;) ALTERNATIVE*; MEMBER*; NOTION LIST.

MC9 MEMBER-

(TERMINALS;) COMMA; OPEN*; CLOSE*; NOTION*.
(NONTERMINALS;) NOTION LIST*.

MC10 NOTION LIST+-

(TERMINALS;) COMMA*; NOTION*.
(NONTERMINALS;) NOTION LIST*.

MAY BE CONTAINED IN (C1): A, B - AB.

C11 COLON, LEFT HAND SIDE-

(NONTERMINALS;) GRAMMAR; RULES; RULE*.

C12 POINT-

(NONTERMINALS;) GRAMMAR; TERMINALS; RULES; REST TERMINALS*; RULE*.

C13 SEMICOLON-

(NONTERMINALS;) GRAMMAR; TERMINALS; RULES; REST TERMINALS*; RULE;
ALTERNATIVES*.

C14 COMMA, NOTION LIST+-

(NONTERMINALS;) GRAMMAR; RULES; RULE; ALTERNATIVES; ALTERNATIVE*;
MEMBER*; NOTION LIST*.

C15 OPEN, CLOSE-

(NONTERMINALS;) GRAMMAR; RULES; RULE; ALTERNATIVES; ALTERNATIVE;
MEMBER*.

C16 NOTION-

(NONTERMINALS;) GRAMMAR; TERMINALS*; RULES; REST TERMINALS; RULE;
LEFT HAND SIDE*; ALTERNATIVES; ALTERNATIVE; MEMBER*;
NOTION LIST*.

C17 EOF-

(NONTERMINALS;) GRAMMAR*.

C18 GRAMMAR-

C19 TERMINALS+, REST TERMINALS+-

(NONTERMINALS;) GRAMMAR*; TERMINALS*; REST TERMINALS*.

C110 RULES+, RULE-

(NONTERMINALS;) GRAMMAR*; RULES*.

C111 ALTERNATIVES+-

(NONTERMINALS;) GRAMMAR; RULES; RULE*; ALTERNATIVES*.

C112 ALTERNATIVE+, MEMBER-

(NONTERMINALS;) GRAMMAR; RULES; RULE; ALTERNATIVES*; ALTERNATIVE*.

MAY BEGIN WITH (BW): AB - A.

BW1 COLON, POINT, SEMICOLON, COMMA, OPEN, CLOSE, NOTION, EOF-

BW2 GRAMMAR-
(TERMINALS:) NOTION.
(NONTERMINALS:) TERMINALS*.

BW3 TERMINALS, LEFT HAND SIDE, NOTION LIST-
(TERMINALS:) NOTION*.

BW4 RULES-
(TERMINALS:) NOTION.
(NONTERMINALS:) RULE*; LEFT HAND SIDE.

BW5 REST TERMINALS-
(TERMINALS:) POINT*; SEMICOLON*.

BW6 RULE-
(TERMINALS:) NOTION.
(NONTERMINALS:) LEFT HAND SIDE*.

BW7 ALTERNATIVES-
(TERMINALS:) SEMICOLON*; OPEN; NOTION.
(NONTERMINALS:) ALTERNATIVE*; MEMBER.

BW8 ALTERNATIVE-
(TERMINALS:) OPEN; NOTION.
(NONTERMINALS:) MEMBER*.

BW9 MEMBER-
(TERMINALS:) OPEN*; NOTION*.

MAY BE THE BEGIN OF (B0): A - AB.

B01 COLON, COMMA, CLOSE, EOF, GRAMMAR, RULES, REST TERMINALS,
ALTERNATIVES, NOTION LIST-

B02 POINT-
(NONTERMINALS:) REST TERMINALS*.

B03 SEMICOLON-
(NONTERMINALS:) REST TERMINALS*; ALTERNATIVES*.

B04 OPEN-
(NONTERMINALS:) ALTERNATIVES; ALTERNATIVE; MEMBER*.

B05 NOTION-
(NONTERMINALS:) GRAMMAR; TERMINALS*; RULES; RULE; LEFT HAND SIDE*;
ALTERNATIVES; ALTERNATIVE; MEMBER*; NOTION LIST*.

B06 TERMINALS-
(NONTERMINALS:) GRAMMAR*.

B07 RULE-
(NONTERMINALS:) RULES*.

B08 LEFT HAND SIDE-
(NONTERMINALS:) RULES; RULE*.

B09 ALTERNATIVE-
(NONTERMINALS:) ALTERNATIVES*.

B010 MEMBER-
(NONTERMINALS:) ALTERNATIVES; ALTERNATIVE*.

MAY END WITH (EW): AB = B.

EW1 COLON, POINT, SEMICOLON, COMMA, OPEN, CLOSE, NOTION, EOF-

EW2 GRAMMAR-
(TERMINALS:) EOF*.

EW3 TERMINALS+, REST TERMINALS+-
(TERMINALS:) POINT*.
(NONTERMINALS:) TERMINALS*; REST TERMINALS*.

EW4 RULES+-
(TERMINALS:) POINT.
(NONTERMINALS:) RULES*; RULE*.

EW5 RULE-
(TERMINALS:) POINT*.

EW6 LEFT HAND SIDE-
(TERMINALS:) NOTION*.

EW7 ALTERNATIVES+-
(TERMINALS:) SEMICOLON*; CLOSE; NOTION.
(NONTERMINALS:) ALTERNATIVES*; ALTERNATIVE*; MEMBER.

EW8 ALTERNATIVE+-
(TERMINALS:) CLOSE; NOTION.
(NONTERMINALS:) ALTERNATIVE*; MEMBER*.

EW9 MEMBER-
(TERMINALS:) CLOSE*; NOTION*.

EW10 NOTION LIST+-
(TERMINALS:) NOTION*.
(NONTERMINALS:) NOTION LIST*.

MAY BE THE END OF (EO): B - AB.

EO1 COLON, COMMA, OPEN, GRAMMAR, LEFT HAND SIDE-

EO2 POINT-

(NONTERMINALS:) TERMINALS; RULES; REST TERMINALS*; RULE*.

EO3 SEMICOLON, ALTERNATIVES+-

(NONTERMINALS:) ALTERNATIVES*.

EO4 CLOSE-

(NONTERMINALS:) ALTERNATIVES; ALTERNATIVE; MEMBER*.

EO5 NOTION-

(NONTERMINALS:) LEFT HAND SIDE*; ALTERNATIVES; ALTERNATIVE; MEMBER*;
NOTION LIST*.

EO6 EOF-

(NONTERMINALS:) GRAMMAR*.

EO7 TERMINALS+, REST TERMINALS+-

(NONTERMINALS:) TERMINALS*; REST TERMINALS*.

EO8 RULES+, RULE-

(NONTERMINALS:) RULES*.

EO9 ALTERNATIVE+, MEMBER-

(NONTERMINALS:) ALTERNATIVES*; ALTERNATIVE*.

EO10 NOTION LIST+-

(NONTERMINALS:) NOTION LIST*.

MAY FOLLOW (MF): B - A.

MF1 COLON-
(TERMINALS:) NOTION.
(NONTERMINALS:) LEFT HAND SIDE*.

MF2 POINT-
(TERMINALS:) COLON*; SEMICOLON; CLOSE; NOTION*.
(NONTERMINALS:) ALTERNATIVES*; ALTERNATIVE; MEMBER.

MF3 SEMICOLON+-
(TERMINALS:) COLON*; SEMICOLON*; CLOSE; NOTION*.
(NONTERMINALS:) ALTERNATIVE*; MEMBER.

MF4 COMMA-
(TERMINALS:) CLOSE; NOTION*.
(NONTERMINALS:) MEMBER*.

MF5 OPEN, ALTERNATIVE, MEMBER-
(TERMINALS:) COLON*; SEMICOLON*; COMMA*.

MF6 CLOSE-
(TERMINALS:) NOTION.
(NONTERMINALS:) NOTION LIST*.

MF7 NOTION-
(TERMINALS:) COLON*; POINT; SEMICOLON*; COMMA*; OPEN*.
(NONTERMINALS:) TERMINALS*; REST TERMINALS; RULE*.

MF8 EOF-
(TERMINALS:) POINT.
(NONTERMINALS:) RULES*; RULE.

MF9 GRAMMAR-

MF10 TERMINALS-
(TERMINALS:) SEMICOLON*.

MF11 RULES, RULE+, LEFT HAND SIDE-
(TERMINALS:) POINT.
(NONTERMINALS:) TERMINALS*; REST TERMINALS; RULE*.

MF12 REST TERMINALS-
(TERMINALS:) NOTION*.

MF13 ALTERNATIVES-
(TERMINALS:) COLON*; SEMICOLON*.

MF14 NOTION LIST-
(TERMINALS:) COMMA*; OPEN*.

MAY PRECEDE (MP): A - B.

MP1 COLON-

(TERMINALS:) POINT*; SEMICOLON; OPEN; NOTION.
(NONTERMINALS:) ALTERNATIVES*; ALTERNATIVE; MEMBER.

MP2 POINT, RULE+-

(TERMINALS:) NOTION; EOF*.
(NONTERMINALS:) RULES*; RULE; LEFT HAND SIDE.

MP3 SEMICOLON+-

(TERMINALS:) POINT*; SEMICOLON; OPEN; NOTION.
(NONTERMINALS:) TERMINALS*; ALTERNATIVES*; ALTERNATIVE; MEMBER.

MP4 COMMA-

(TERMINALS:) OPEN; NOTION.
(NONTERMINALS:) ALTERNATIVE*; MEMBER; NOTION LIST*.

MP5 OPEN-

(TERMINALS:) NOTION.
(NONTERMINALS:) NOTION LIST*.

MP6 CLOSE, MEMBER-

(TERMINALS:) POINT*; SEMICOLON*; COMMA*.

MP7 NOTION-

(TERMINALS:) COLON*; POINT*; SEMICOLON*; COMMA*; CLOSE*.
(NONTERMINALS:) REST TERMINALS*.

MP8 EOF, GRAMMAR-

MP9 TERMINALS, REST TERMINALS-

(TERMINALS:) NOTION.
(NONTERMINALS:) RULES*; RULE; LEFT HAND SIDE.

MP10 RULES-

(TERMINALS:) EOF*.

MP11 LEFT HAND SIDE-

(TERMINALS:) COLON*.

MP12 ALTERNATIVES-

(TERMINALS:) POINT*.

MP13 ALTERNATIVE-

(TERMINALS:) POINT*; SEMICOLON*.

MP14 NOTION LIST-

(TERMINALS:) CLOSE*.

ALTERNATIVE
 ALTERNATIVES
 CLOSE
 COLON
 COMMA
 EOF
 GRAMMAR
 LEFT HAND SIDE
 MEMBER
 NOTION
 NOTION LIST
 OPEN
 POINT
 REST TERMINALS
 RULE
 RULES
 SEMICOLON
 TERMINALS

MC8, C12, BW8, BO9, EW8, EO9, MF0, MP13
 MC7, C11, BW7, BO0, EW7, EO0, MF13, MP12
 MC0, C11, BW5, BO1, EW1, EO4, MF6, MP6
 MC1, C11, BW1, BO1, EW1, EO1, MF1, MP1
 MC0, C14, BW0, BO1, EW1, EO1, MF4, MP4
 MC0, C17, BW0, BO1, EW1, EO6, MF8, MP8
 MC2, C18, BW2, BO0, EW2, EO0, MF9, MP0
 MC6, C10, BW1, BO8, EW6, EO0, MF1, MP11
 MC9, C10, BW9, BO10, EW9, EO0, MF9, MP5
 MC0, C16, BW0, BO5, EW0, EO5, MF7, MP7
 MC10, C10, BW4, BO3, EW10, EO10, MF14, MP14
 MC0, C15, BW0, BO4, EW0, EO1, MF5, MP5
 MC0, C12, BW0, BO2, EW0, EO2, MF2, MP2
 MC0, C13, BW5, BO0, EW1, EO3, MF12, MP0
 MC5, C10, BW6, BO7, EW5, EO0, MF8, MP11
 MC4, C10, BW4, BO0, EW4, EO8, MF11, MP10
 MC0, C13, BW0, BO3, EW0, EO3, MF3, MP3
 MC3, C19, BW3, BO6, EW3, EO7, MF10, MP9

3. LL(1)-checker

The necessary requirements for a grammar, written in a form without optional parts and without useless non-terminals (i.e., without non-terminals which either do not produce any finite string, or do not depend on the root of the grammar), to be of type LL(1), are:

1. for any rule of the form $A:a_1;\dots;a_n$, the sets $\text{first}(a_1), \dots, \text{first}(a_n)$, where $\text{first}(a_i) = \{s | a_i \text{ FIRST}^*s\}$, are mutually disjoint;
2. at most one of the $a_1, \dots, a_n$ can produce the null string (ϵ);
3. if a_p produces ϵ , then $\text{first}(A)$ has no elements in common with $\text{follow}(A)$, where $\text{follow}(A) = \{s | s \text{ "may follow" } A\}$.

This third requirement is slightly different from, but equivalent to, the one given by Knuth [3], and somewhat easier to check. The equivalence is proved in [1].

The check for LL(1)-ness is done in several steps, corresponding to the requirements listed above:

- i. Requirement 1 is checked in two steps: during the computation of FIRST it is tested that no elements of the array "first" are filled twice, i.e., that there are no non-terminals for which two alternatives "directly start" with the same notion (either terminal or non-terminal). Secondly, all "direct starts", i.e., notions that are the beginning of some non-terminal, are looked at to test that there are no direct starts that may (directly or indirectly) begin with the same terminal. (The test for non-terminals is, obviously, superfluous here.)
- ii. Requirement 2 is also checked in two steps: first, it is tested that non-terminals that produce empty (ϵ), do not have more than one alternative producing ϵ (which is exactly the requirement as stated before). However, as our grammar has optional parts, the full test is slightly more complicated. If each optional part is replaced by some non-terminal, the production rule for which has two alternatives: ϵ , and the enclosed list of notions, requirement 2 is fulfilled if the enclosed list

of notions does not produce ϵ . This is tested along with requirement 3. iii. Requirement 3 is checked more or less similarly to requirement 2. The check is done for notions that produce ϵ , and for the implicit notions for which the enclosed parts stand.

The test for useless non-terminals is not done. Notions which do not depend on the root of the grammar are of no importance at all, notions which do not produce any finite string are not taken into account.

In section 3.1 a listing of the program is given, section 3.2 shows two small examples, one for a grammar which is not of type $LL(1)$, and one for a grammar which is $LL(1)$ and generates the same language.

3.1 The program

begin comment ll(1)-checker, august-september 1973;

integer end of file, eof, space char, tab char, nlcr char,
 sub char, bus char, stock, nil, comma char, comma,
 point char, point, colon char, colon, semicolon char,
 semicolon, open char, open, close char, close,
 list pointer, max int, first notion, new defined,
 head of tree, top, error count, sym, number, numb,
 notion, max notion, start notion, harmless, minus char,
 n of adm cells, upperbound, star, symbols, letter a,
 space repr, bits per word;
boolean ll1;
integer array t32[0 : 4];

initialization part:

end of file:= eof:= -4096; space char:= 93; tab char:= 118;
 nlcr char:= 119; sub char:= 100; bus char:= 101; nil:= -1;
 comma char:= comma:= 87; point char:= 88; point:= -5;
 colon char:= colon:= 90; semicolon char:= 91; semicolon:= -4;
 open char:= 98; open:= -2; close char:= 99; close:= -3;
 max int:= 67 108 863; first notion:= 1; new defined:= -1;
 t32[0]:= 1 048 576; t32[1]:= 32 768; t32[2]:= 1 024; t32[3]:= 32;
 t32[4]:= 1; list pointer:= 1; star:= 66; error count:= 0;
 harmless:= -1; n of adm cells:= 3;
 minus char:= 65; letter a:= 10; space repr:= 27; bits per word:= 27;
 max notion:= 300;
 upperbound:= top:= available - max notion - 5 × max notion ×
 max notion / bits per word - 2000;
 printtext(↑upperbound array list:↓); absfixt(5, 0, upperbound);
 nlcr; nlcr;

begin integer array list[1 : upperbound],
 notion link[1 : max notion];
comment 'notion link' contains pointers to notions, stored in
 'list', as follows:
 left
 right
 number
 pointer: text 1
 :
 text last - max int,
 starting from 1 up to 'listpointer'.
 the essential part of the grammar is stored in 'list'
 from 'upperbound' to 'top' (i.e. backwards);
boolean array first, first star, last, last star, follow
 [1 : max notion, 1 : max notion],
 possibly empty[1 : max notion];

```

comment reading department;

integer procedure char;
begin integer i, j;
repeat: i:= char:= resymbol;
  if i ≠ end of file then prsym(i);
  if i = tab char v i = nlcr char then goto repeat else
  if i = space char then
  begin j:= resymbol;
    for i:= j while i = space char v i = tab char v
    i = nlcr char v i = sub char do
    begin if i = sub char then
      begin for i:= sub char, j while j ≠ bus char ^
      j ≠ end of file do
        begin prsym(i); j:= resymbol end;
        if j = bus char then
        begin prsym(j); j:= resymbol end
        else error(⌘premature end of file⌘)
      end
    else
      begin prsym(i); j:= resymbol end
    end;
    if j ≠ end of file then prsym(j);
    if letter(j) then stock:= j else char:= j
  end
  else if i = sub char then
skip comment:
  begin i:= char:= resymbol; prsym(i);
    if i = bus char then goto repeat;
    if i ≠ end of file then goto skip comment
  end
end char;

procedure next symbol;
comment this procedure yields a symbol in "sym", and,
  if it is a notion, yields its number in "numb";
begin
again: if sym = end of file then else
  if stock = nil then sym:= char else
  begin sym:= stock; stock:= nil end;
  if sym = space char then goto again else
  if sym = comma char then sym:= comma else
  if sym = point char then sym:= point else
  if sym = colon char then sym:= colon else
  if sym = semicolon char then sym:= semicolon else
  if sym = open char then sym:= open else
  if sym = close char then sym:= close else
  if sym = end of file then sym:= eof else
  if letter(sym) then
  begin integer j, aux, sm;
    number:= number + 1;

```

```

for j:= nil, nil, number do store(j);
aux:= j:= 0; start notion:= list pointer;
for sm:= sym, sm while letter(sm) do
begin if j = 5 then
begin store(aux); aux:= j:= 0 end store in word;
aux:= t32[j] × (if sm = space char then space repr
else sm - letter a + 1) + aux; j:= j + 1;
if stock = nil then sm:= char else
begin sm:= stock; stock:= nil end
end pack symbols;
store( - max int + aux); stock:= sm;
numb:= in tree(start notion, head of tree);
if numb = new defined then
begin numb:= new(number); notion link[number]:= - start
notion
end not yet defined
else
if defined(numb) then numb:= new(numb);
sym:= notion
end read and store a notion
else
begin error({illegal character; skipped}); goto again end
end next symbol;

```

```

integer procedure in tree(start, last ref);
value start, last ref; integer start, last ref;
comment this procedure considers the notion at "start". if this
notion is in the tree, it yields its number and removes
the notion, otherwise, it adds the notion to the tree
and yields "new defined";
begin integer node, comp;
boolean found;

```

```

integer procedure compare(one, two); value one, two;
integer one, two;
comment compare:= sign(one '-' two);
begin integer o, t, so, st, i;
o:= list[one]; t:= list[two];
for i:= i while o = t ^ o > 0 do
begin one:= one + 1; o:= list[one]; two:= two + 1;
t:= list[two]
end;
comment discriminating item found, now analyze it;
if o < 0 then
begin o:= o + max int; so:= - 1 end
else so:= 1;
if t < 0 then
begin t:= t + max int; st:= - 1 end
else st:= 1;
compare:= sign(if o = t then so - st else o - t)
end compare;

```

```

    found:= false;
    comment search tree;
    for node:= list[last ref] while node ≠ nil ^ ¬ found do
    begin comp:= compare(node, start);
        if comp = 0 then found:= true else
            last ref:= node - n of adm cells +
                (if comp < 0 then 1 else 0)
    end node;
    comment analyze result: ;
    if found then
    begin listpointer:= start - n of adm cells;
        number:= number - 1; comment item removed;
        in tree:= list[node - 1]
    end list[head - 1] contains the number of the notion
    else
    begin list[last ref]:= start; in tree:= new defined
    end item added
end in tree;

boolean procedure letter(sym); value sym; integer sym;
letter:= letter a ≤ sym ^ sym < letter a + 26 v sym = space char;

procedure store(sym); value sym; integer sym;
if list pointer > top then
begin carriage(2); printtext(␣space exhausted␣); exit end
else
begin list[list pointer]:= sym; list pointer:= list pointer + 1
end store;

procedure store on top(sym); value sym; integer sym;
if top < list pointer then
begin carriage(2); printtext(␣space exhausted␣); exit end
else
begin list[top]:= sym; top:= top - 1 end store on top;

integer procedure new(number); value number; integer number;
new:= - number;

boolean procedure is new(number); value number;
integer number;
is new:= number < 0;

procedure define(number); value number; integer number;
notion link[number]:= abs(notion link[number]);

boolean procedure defined(number); value number;
integer number;
defined:= notion link[number] > 0;

```

comment the syntax of the input is given by:

```
grammar: terminals, rules, eof.
```

```
terminals: notion, point|
          notion, semicolon, terminals.
```

```
rules: rule, (rules).
```

```
rule: notion, colon, tail, point.
```

```
tail: (alternative), (semicolon, tail).
```

```
alternative: member, (comma, member).
```

```
member: notion|
```

```
        open, notion list, close.
```

```
notion list: notion, (comma, notion list).
```

the next part reads in such a grammar, puts its structure in the back end of the array "list", and puts the notions in a tree, stored at the front end of "list". obviously, the | stands for ;

```
procedure grammar;
```

```
begin number:= 0; sym:= stock:= nil; head of tree:= list pointer;
```

```
        store(nil); next symbol; terminals; rules
```

```
end grammar;
```

```
procedure terminals;
```

```
terminals:
```

```
begin req notion(false);
```

```
        if is new(num) then define(new(num)) else
```

```
        error({terminal occurs twice});
```

```
        if is(semicolon, false) then goto terminals else
```

```
        if is(point, false) then symbols:= number - 1 else
```

```
        if is(eof, false) then else
```

```
        begin error({error in terminals}); next symbol;
```

```
        goto terminals
```

```
        end
```

```
end terminals;
```

```

procedure rules;
rules:
begin req notion(true);
    if is new(num) then define(new(num)) else
        error({notion already defined});
    if is(colon, false) then else error({colon missing});
tail: if is(point, true) then
    begin if is(eof, false) then else goto rules end
    else if is(semicol, true) then goto tail else
        if member then
rest list:
    begin if is(comma, false) then
        begin if ¬ member then error({alternative wrong});
            goto rest list
        end
        else goto tail
    end
    else if is(eof, false) then
        begin store on top(point); error({premature end of file});
        end
    else
        begin error({incorrect rule}); next symbol; goto tail end
    end rules;

boolean procedure member;
if is(notion, true) then member:= true else
if is(open, true) then
begin member:= true;
rest member: req notion(true);
    if is(comma, false) then goto rest member else
    if is(close, true) then else error({option wrong})
end
else member:= false;

procedure req notion(copy); value copy; boolean copy;
if ¬ is(notion, copy) then
begin error({notion missing}); numb:= harmless end req notion;

boolean procedure is(s, copy); value s, copy; integer s;
boolean copy;
if s = sym then
begin is:= true; if copy then
    store on top(if sym = notion then abs(num) else sym);
    next symbol
end
else is:= false;

procedure error(s); string s;
begin integer pos;
    pos:= printpos; error count:= error count + 1; carriage(2);
    printtext({error: }); printtext(s); carriage(3); space(pos)
end error;

```

```

comment compute "possibly empty", "first", "first star", "last",
    "last star" and "follow";

procedure check for empty productions;
comment check which notions produce empty. the procedure continues
    until there are no more changes;
begin integer aux, notion, el, i;
    boolean changed, any;
    any := false;
    for aux := 1 step 1 until number do
        possibly empty[aux] := false;
    check for empty productions: changed := false;
    for aux := upperbound, aux - 1 while aux > top do
        begin notion := list[aux];
            if possibly empty[notion] then
                begin for el := list[aux] while el ≠ point do
                    aux := aux - 1
                end was already empty
            else
                begin
                    next: aux := aux - 1; el := list[aux]; if el = open then
                        begin for el := list[aux] while el ≠ close do
                            aux := aux - 1; goto next
                        end optional part
                    else if el = semicolon v el = point then
                        begin for el := list[aux] while el ≠ point do
                            aux := aux - 1;
                            possibly empty[notion] := changed := any := true
                        end empty production found
                    else if possibly empty[el] then goto next else
                        for el := list[aux] while el ≠ point do
                            if el = semicolon then goto next else
                                aux := aux - 1
                            end production rule
                        end grammar;
                    if changed then goto check for empty productions;
                    if any then
                        begin printtext({the following notions may produce empty:});
                            nlcr;
                            for i := 1 step 1 until number do
                                if possibly empty[i] then
                                    begin nlcr; print notion(i) end
                            end
                        else printtext({no rule produces empty});
                            carriage(5);
                            check only one alternative yields empty
                        end check for empty productions;
                    end
                end
            end
        end
    end

```

```

procedure check only one alternative yields empty;
comment this procedure checks whether not more than one
        alternative of a possibly empty notion yields "empty";
begin integer aux, notion, el, numb of empties;
        for aux:= upperbound, aux - 1 while aux > top do
            begin notion:= list[aux];
                if possibly empty[notion] then
                    begin numb of empties:= 0;
                        next: aux:= aux - 1; el:= list[aux];
                            if el = open then
                                begin for el:= list[aux] while el ≠ close do
                                    aux:= aux - 1; goto next
                                end optional part
                            else if el = semicolon ∨ el = point then
                                begin numb of empties:= numb of empties + 1;
                                    if numb of empties > 1 then
                                        begin ll1:= false; message1(⊥in ⊥, notion,
                                            ⊥ two alternatives yield empty⊥)
                                        end;
                                        if el = semicolon then goto next
                                    end alternative
                                else if possibly empty[el] then goto next else
                                    for el:= list[aux] while el ≠ point do
                                        if el = semicolon then goto next else
                                            aux:= aux - 1
                                        end empty notion
                                    else
                                        for el:= list[aux] while el ≠ point do aux:= aux-1
                                    end grammar
                                end check only one alternative yields empty;

```

```

procedure may begin with;
comment this procedure determines the relation "may begin with".
    the result is stored in 'first' (directly beginning with),
    its transitive closure in 'first star'.
    while determining 'first', error messages are given for
    the "direct initial uncertainties" found;
begin integer i, j, el, aux, notion;
    boolean optional, errors, any;
    errors := false;
    for i := 1 step 1 until number do
    for j := 1 step 1 until number do
    first[i, j] := first star[i, j] := false;
    for aux := upperbound, aux - 1 while aux > top do
    begin notion := list[aux]; aux := aux - 1; optional := false;
        for el := list[aux] while el ≠ point do
        begin aux := aux - 1;
            if el = open then optional := true else
            if el = close then optional := false else
            if el ≠ semicolon then
            begin if first[notion, el] then
                begin errors := true;
                    message2({two alternatives in †, notion,
                        † start with †, el})
                end direct uncertainty
            else first[notion, el] := first star[notion, el] :=
                true;
                if possibly empty[el] then else
                if optional then
                begin for el := list[aux] while el ≠ close
                    do aux := aux - 1
                end skip rest of option
            else
                for el := list[aux] while el ≠ semicolon ^
                el ≠ point do aux := aux - 1
            end
        end production rule
    end grammar;
    if errors then ll1 := false else
    printtext({no direct initial uncertainties found}); carriage(5);
    transitive closure(first star);
    any := false;
    for i := 1 step 1 until number do
    if first star[i, i] then
    begin if ¬ any then
        begin printtext({the following rules are left-recursive:});
            nlcr; any := true
        end;
        nlcr; print notion(i)
    end;
    if ¬ any then printtext({no rule is left-recursive});
    carriage(5);
end may begin with;

```

```

procedure may end with;
comment this procedure determines the relation "may end with";
begin integer i, j, n, p, aux, notion, el;
    boolean optional;
    for i:= 1 step 1 until number do
    for j:= 1 step 1 until number do
    last[i, j]:= last star[i, j]:= false;
    for aux:= upperbound, aux while aux > top do
    begin notion:= list[aux]; n:= aux; optional:= false;
        for el:= list[aux] while el ≠ point, el do
        aux:= aux - 1; p:= aux + 2;
        for el:= list[p] while p < n do
        begin p:= p + 1;
            if el = open then optional:= false else
            if el = close then optional:= true else
            if el ≠ semicolon then
            begin last[notion, el]:= last star[notion, el]:=
                true;
                if possibly empty[notion] then else
                if optional then
                begin for el:= list[p] while el ≠ open do
                    p:= p + 1
                end skip rest of option
                else
                for el:= list[p] while el ≠ semicolon ^ p < n
                do p:= p + 1
            end
        end production rule
    end grammar;
    transitive closure(last star)
end may end with;

procedure transitive closure(r); boolean array r;
comment warshall algorithm, see "grammar handling tools", p48;
begin integer i, j, k;
    for i:= 1 step 1 until number do
    for j:= symbols + 1 step 1 until number do
    if r[j, i] then
    begin for k:= 1 step 1 until number do
        if r[i, k] then r[j, k]:= true
    end
end transitive closure;

```

```

procedure follow within;
comment this procedure determines the successions of notions
        within the production rules;
begin integer i, j, stackpointer, lwb stack, upb stack, el,
        previous, aux;
        integer array stack[1 : 100];
        comment the upperbound 100 may be erroneous, a good upperbound
                is the maximum number of notions in an alternative,
                but this will very likely be less than 100;
        for i:= 1 step 1 until number do
        for j:= 1 step 1 until number do follow[i, j]:= false;
        for aux:= upperbound, aux - 1 while aux > top do
        for el:= semicolon, list[aux] while el ≠ point do
        begin aux:= aux - 1;
            if el = open then
                begin if previous ≠ nil then
                    begin stackpointer:= stackpointer + 1;
                        stack[stackpointer]:= previous
                    end;
                    upb stack:= stackpointer + 1
                end
            else if el = close then lwb stack:= upb stack:= 1
            else if el = semicolon then
                begin stackpointer:= 0; lwb stack:= upb stack:= 1;
                    previous:= nil
                end
            else
                begin for i:= lwb stack step 1 until stackpointer
                    do follow[stack[i], el]:= true;
                    if previous ≠ nil then
                        begin follow[previous, el]:= true;
                            if possibly empty[el] then
                                begin stackpointer:= stackpointer + 1;
                                    stack[stackpointer]:= previous
                                end
                            else
                                begin stackpointer:= upb stack - 1;
                                    lwb stack:= upb stack
                                end
                            end;
                            previous:= el
                        end notion
                    end rule and grammar
                end follow within;

```

```

comment check for ll(1)-ness;

procedure report indirect initial uncertainties;
comment this procedure determines whether two alternatives of
    one rule start with the same terminal symbol;
begin integer n1;
    boolean any, left;

    procedure report(n1, n2); value n1, n2; integer n1, n2;
    begin if  $\neg$  any then
        begin printtext(4for the following notions, more than 4);
            printtext(4one alternative may4); nclr; space(5);
            printtext(4start with a given notion:4); nclr;
            any := true; ll1 := false
        end;
        nclr; if left then
            begin nclr; print notion(n1); prsym(minus char); nclr;
            left := false
        end;
        space(8); print notion(n2)
    end report;

    any := false;
    for n1 := symbols + 1 step 1 until number do
        begin integer count, n2, i, si, j, sj;
            integer array direct start[1 : number];
            comment collect "direct start"s of n1;
            left := true; count := 0;
            for n2 := 1 step 1 until number do
                if first[n1, n2] then
                    begin count := count + 1; direct start[count] := n2 end;
                    for i := 1 step 1 until count do
                        begin si := direct start[i];
                            for j := i + 1 step 1 until count do
                                begin sj := direct start[j];
                                    if si < symbols then
                                        begin if sj < symbols then else
                                            if first star[sj, si] then report(n1, si)
                                        end si is basic, sj may start with si
                                        else
                                            for n2 := 1 step 1 until symbols do
                                                if first star[si, n2]  $\wedge$  first star[sj, n2] then
                                                    report(n1, n2)
                                            end sj
                                        end si
                                    end rules;
                                if  $\neg$  any then
                                    printtext(4no indirect initial uncertainties found4);
                                    carriage(5)
                                end report indirect initial uncertainties;
                            end sj
                        end i
                    end j
                end n2
            end count
        end n1
    end n2

```

```

procedure report indirect uncertainties;
comment this procedure detects violations of requirement 3;
begin integer aux, el, handle;
      boolean optional, any;

  procedure check follow(start, in first option);
  value start, in first option; integer start;
  boolean in first option;
  comment checks whether the notion "i" at "start" (which is
           possibly empty or the beginning of an optional part)
           may be followed(in the rule for "handle") by a notion
           which starts with a terminal symbol that may also be
           the beginning of "i";
  begin integer aux, el;
    boolean test completed, in next option;
    aux:= start - 1; start:= list[start];
    in next option:= test completed:= false;
    for el:= list[aux] while  $\neg$  test completed do
      begin if el = open then in next option:= true else
        if el = close then
          in next option:= in first option:= false else
            if el = semicolon v el = point then
              begin check end(start); test completed:= true end
            else
              begin if possibly empty[start] v  $\neg$  in first option
                then
                  begin integer i;
                    for i:= 1 step 1 until symbols do
                      if first star[start, i]  $\wedge$  first star[el, i]
                        then
                          begin any:= true; message6( $\{$ in  $\{$ , handle,
                             $\{$  the possibly empty or optional notion  $\{$ ,
                            start,  $\{$ may be followed by  $\{$ , el,  $\{$ ;  $\{$ ,
                             $\{$ both  $\{$ , start,  $\{$  and  $\{$ , el,
                             $\{$  may begin with  $\{$ , i $\}$ ); ll1:= false
                          end message
                        end compare the beginning of start and el;
                        if possibly empty[el] then else
                          if in first option v in next option then
                            begin for el:= list[aux - 1] while el  $\neq$ 
                              close do aux:= aux - 1
                            end skip rest of option
                          else test completed:= true
                        end notion;
                        aux:= aux - 1
                      end test
                    end check follow;

```

```

procedure check end(start); value start; integer start;
comment similar to "check follow", but now "i" is the
beginning of an optional last member or is a
possibly empty last notion of "handle";
begin integer i, j;
  for i:= symbols + 1 step 1 until number do
    if last star[i, handle] then
      begin comment i may end with handle;
        for j:= 1 step 1 until number do
          if follow[i, j] then
            begin comment j may follow i;
              integer k;
              for k:= 1 step 1 until symbols do
                if first star[start, k]  $\wedge$  first star[j, k] then
                  begin message7({the notion  $\downarrow$ , start,
                     $\downarrow$  is the beginning of the possibly empty $\downarrow$ ,
                     $\downarrow$  or optional last member of  $\downarrow$ , i,  $\downarrow$  (via  $\downarrow$ ,
                    handle,  $\downarrow$ ) $\downarrow$ ,  $\downarrow$  and may be followed by  $\downarrow$ , j,  $\downarrow$ ;  $\downarrow$ ,
                     $\downarrow$  both  $\downarrow$ , start,  $\downarrow$  and  $\downarrow$ , j,  $\downarrow$  may begin with  $\downarrow$ ,
                    k}); any:= ll1:= false
                  end message
                end compare the beginning of start and j
              end
            end
          end
        end
      end
    end
  check end;

procedure check against ambiguous option(start);
value start; integer start;
comment checks whether an option starting at "start" does not
contain only empty-producing notions, i.e. is
ambiguous;
begin integer aux;
  boolean ok;
  for aux:= start - 1, aux - 1 while ok do
    begin el:= list[aux]; if el = close then
      begin ok:= ll1:= false; any:= true; message3({in  $\downarrow$ ,
        handle,  $\downarrow$  the optional part starting with  $\downarrow$ ,
        start - 1,  $\downarrow$  produces empty in more than one way $\downarrow$ )
      end
      else ok:= possibly empty[el]
    end
  end option
end check against ambiguous option;

for el:= 1 step 1 until number do
  first star[el, el]:= last star[el, el]:= true;
  printtext({violations of requirement 3 $\downarrow$ }); nlcr;
  printtext({(i.e., ambiguities arising from empty notions or  $\downarrow$ )});
  printtext({optional parts: $\downarrow$ }); nlcr; nlcr; any:= false;
  for aux:= upperbound, aux - 1 while aux > top do
    begin handle:= list[aux]; aux:= aux - 1; optional:= false;

```

```

    for el:= list[aux] while el ≠ point do
    begin if el = open then
        begin optional:= true;
            check against ambiguous option(aux)
        end
        else
            if el = close then optional:= false else
            if el ≠ semicolon then
                begin if possibly empty[el] then
                    check follow(aux, optional) else
                    if optional then
                        begin check follow(aux, true); optional:=false
                        end first [not empty] notion of optional part
                    end;
                end;
                aux:= aux - 1
            end rule
        end grammar;
        if ¬ any then printtext(⟨none⟩)
    end report indirect uncertainties;

comment output department;

procedure message1(s1, n1, s2); value n1; integer n1;
string s1, s2;
begin nlcr; printtext(s1); print notion(n1);
    printtext(s2); nlcr
end message1;

procedure message2(s1, n1, s2, n2); value n1, n2;
integer n1, n2; string s1, s2;
begin nlcr; printtext(s1); print notion(n1);
    printtext(s2); print notion(n2); nlcr
end message2;

procedure message3(s1, n1, s2, n2, s3); value n1, n2;
integer n1, n2; string s1, s2, s3;
begin nlcr; printtext(s1); print notion(n1); printtext(s2);
    print notion(n2); printtext(s3); nlcr
end message3;

procedure message6(s1, n1, s2, n2, s3, n3, s4, s5, n4, s6, n5, s7,
    n6);
value n1, n2, n3, n4, n5, n6; integer n1, n2, n3, n4, n5, n6;
string s1, s2, s3, s4, s5, s6, s7;
begin nlcr; nlcr; printtext(s1); print notion(n1);
    printtext(s2); print notion(n2); nlcr; space(5); printtext(s3);
    print notion(n3); printtext(s4); nlcr; space(5); printtext(s5);
    print notion(n4); printtext(s6); print notion(n5);
    printtext(s7); print notion(n6)
end message6;

```

```

procedure message7(s1, n1, s2, s3, n2, s4, n3, s5, s6, n4, s7, s8,
  n5, s9, n6, s10, n7);
value n1, n2, n3, n4, n5, n6, n7; integer n1, n2, n3, n4, n5,
n6, n7; string s1, s2, s3, s4, s5, s6, s7, s8, s9, s10;
begin nlcr; nlcr; printtext(s1); print notion(n1); printtext(s2);
  nlcr; space(5); printtext(s3); print notion(n2);
  nlcr; space(5); printtext(s4); print notion(n3); printtext(s5);
  nlcr; space(5); printtext(s6); print notion(n4); printtext(s7);
  nlcr; space(5); printtext(s8); print notion(n5); printtext(s9);
  print notion(n6); printtext(s10); print notion(n7)
end message7;

```

```

procedure print list of notions;
begin integer i;
  printtext({list of notions, each notion preceded by its number});
  nlcr; printtext({and, if not defined by an asterisk}); nlcr;
  for i:= 1 step 1 until number do
    begin nlcr; if  $\neg$  defined(i) then
      begin prsym(star); define(i); error count:= error count+1
      end
      else prsym(spacechar);
      absfixt(4, 0, i); print notion(i)
    end
  end print list of notions;

```

```

procedure print notion(number); value number; integer number;
begin integer i, n;
  integer array a[1 : 100];
  comment this upperbound is rather arbitrary, notions of
    length > 100 are considered rare;
  get notion at(notion link[number]) of length:(n) in:(a);
  for i:= 1 step 1 until n do prsym(a[i])
end print notion;

```

```

procedure get notion at(index) of length:(n) in:(a);
value index; integer index, n; integer array a;
begin integer el, j, k;
  n:= 0;
  for el:= list[index] while el > 0, el + max int do
    begin for k:= 0, k + 1 while k < 5 ^ el > 0 do
      begin j:= el div t32[k]; el:= el - j x t32[k]; n:= n +1;
      a[n]:= if j = space repr then space char else
        j + letter a - 1
      end;
      index:= index + 1
    end
  end
end get notion at;

```

main program:

```

  grammar;
  newpage;
  print list of notions;
  if error count > 0 then
  begin carriage(3); absfixt(3, 0, error count);
    printtext(4 errors in grammar4); exit
  end;
  newpage;
  ll1:= true;
  check for empty productions;
  may begin with;
  report indirect initial uncertainties;
  may end with;
  follow within;
  report indirect uncertainties;
  carriage(5); printtext(4the grammar is 4);
  if ¬ ll1 then printtext(4not 4); printtext(4of type ll(1)4)
end
end

```

3.2 Two examples

UPPERBOUND ARRAY LIST: 20452

[INPUT:]

[TERMINAL SYMBOLS]

COLON; POINT; SEMICOLON; COMMA; OPEN; CLOSE; NOTION; EOF.

GRAMMAR: TERMINALS, RULES, EOF.

TERMINALS: NOTION, POINT; NOTION, SEMICOLON, TERMINALS.

RULES: RULE; RULE, RULES.

RULE: LEFT HAND SIDE, COLON, ALTERNATIVES, POINT.

LEFT HAND SIDE: NOTION.

ALTERNATIVES: ALTERNATIVE OPTION;

ALTERNATIVE OPTION, SEMICOLON, ALTERNATIVES.

ALTERNATIVE OPTION: ALTERNATIVE; .

ALTERNATIVE: MEMBER; MEMBER, COMMA, ALTERNATIVE.

MEMBER: NOTION; OPEN, NOTION LIST, CLOSE.

NOTION LIST: NOTION; NOTION, COMMA, NOTION LIST.

THE FOLLOWING NOTIONS MAY PRODUCE EMPTY:

ALTERNATIVES
ALTERNATIVE OPTION

TWO ALTERNATIVES IN TERMINALS START WITH NOTION

TWO ALTERNATIVES IN RULES START WITH RULE

TWO ALTERNATIVES IN ALTERNATIVES START WITH ALTERNATIVE OPTION

TWO ALTERNATIVES IN ALTERNATIVE START WITH MEMBER

TWO ALTERNATIVES IN NOTION LIST START WITH NOTION

NO RULE IS LEFT-RECURSIVE

NO INDIRECT INITIAL UNCERTAINTIES FOUND

VIOLATIONS OF REQUIREMENT 3
(I.E., AMBIGUITIES ARISING FROM EMPTY NOTIONS OR OPTIONAL PARTS):

NONE

THE GRAMMAR IS NOT OF TYPE LL(1)

UPPERBOUND ARRAY LIST: 20452

[INPUT:]

[TERMINAL SYMBOLS]

COLON; POINT; SEMICOLON; COMMA; OPEN; CLOSE; NOTION; EOF.

GRAMMAR: TERMINALS, RULES, EOF.

TERMINALS: NOTION, REST TERMINALS.

REST TERMINALS: POINT; SEMICOLON, TERMINALS.

RULES: RULE, (RULES).

RULE: LEFT HAND SIDE, COLON, ALTERNATIVES, POINT.

LEFT HAND SIDE: NOTION.

ALTERNATIVES: (ALTERNATIVE), (SEMICOLON, ALTERNATIVES).

ALTERNATIVE: MEMBER, (COMMA, ALTERNATIVE).

MEMBER: NOTION; OPEN, NOTION LIST, CLOSE.

NOTION LIST: NOTION, (COMMA, NOTION LIST).

THE FOLLOWING NOTIONS MAY PRODUCE EMPTY:
ALTERNATIVES

NO DIRECT INITIAL UNCERTAINTIES FOUND

NO RULE IS LEFT-RECURSIVE

NO INDIRECT INITIAL UNCERTAINTIES FOUND

VIOLATIONS OF REQUIREMENT 3
(I.E., AMBIGUITIES ARISING FROM EMPTY NOTIONS OR OPTIONAL PARTS):

NONE

THE GRAMMAR IS OF TYPE LL(1)

References

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- [3] D.E. Knuth, "Top-down syntax analysis", Acta Informatica 1 (1971), p. 79-110, Springer-Verlag, Berlin.

